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INTRODUCTION

- What Operating Systems Do
- Computer-System Organization
- Computer-System Architecture
- Operating-System Structure
- Operating-System Operations
- Process Management
- Memory Management
- Storage Management
- Protection and Security
- Kernel Data Structures
- Computing Environments
- Open-Source Operating Systems

OBJECTIVES

- To describe the basic organization of computer systems
- To provide a grand tour of the major components of operating systems
- To give an overview of the many types of computing environments
- To explore several open-source operating systems

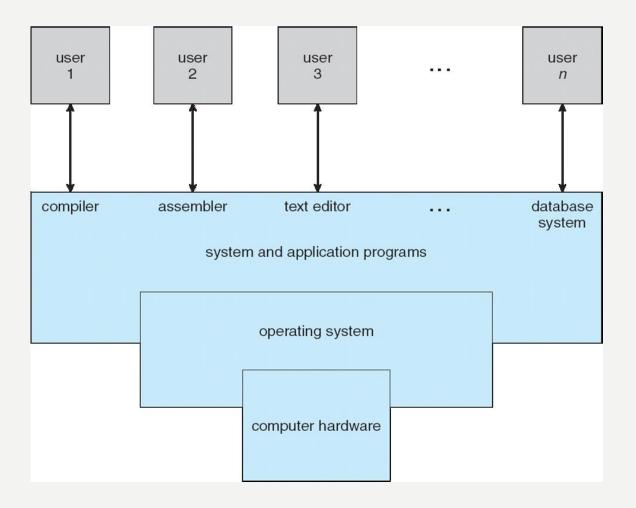
WHAT IS AN OPERATING System?

- A program that acts as an intermediary between a user of a computer and the computer hardware
- Operating system goals:
 - Execute user programs and make solving user problems easier
 - Make the computer system convenient to use
 - Use the computer hardware in an efficient manner

COMPUTER SYSTEM STRUCTURE

- Computer system can be divided into four components:
 - Hardware provides basic computing resources
 - CPU, memory, I/O devices
 - Operating system
 - Controls and coordinates use of hardware among various applications and users
 - Application programs define the ways in which the system resources are used to solve the computing problems of the users
 - Word processors, compilers, web browsers, database systems, video games
 - Users
 - People, machines, other computers

FOUR COMPONENTS OF A COMPUTER SYSTEM

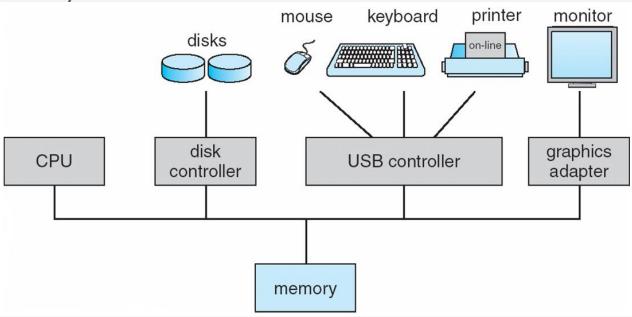


COMPUTER STARTUP

- bootstrap program is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, generally known as firmware
 - Initializes all aspects of system
 - Loads operating system kernel and starts execution

COMPUTER SYSTEM ORGANIZATION

- Computer-system operation
 - One or more CPUs, device controllers connect through common bus providing access to shared memory
 - Concurrent execution of CPUs and devices competing for memory cycles



I/O STRUCTURE

- After I/O starts, control returns to user program only upon I/O completion
 - Wait instruction idles the CPU until the next interrupt
 - Wait loop (contention for memory access)
 - At most one I/O request is outstanding at a time, no simultaneous I/O processing
- After I/O starts, control returns to user program without waiting for I/O completion
 - System call request to the OS to allow user to wait for I/O completion
 - Device-status table contains entry for each I/O device indicating its type, address, and state
 - OS indexes into I/O device table to determine device status and to modify table entry to include interrupt

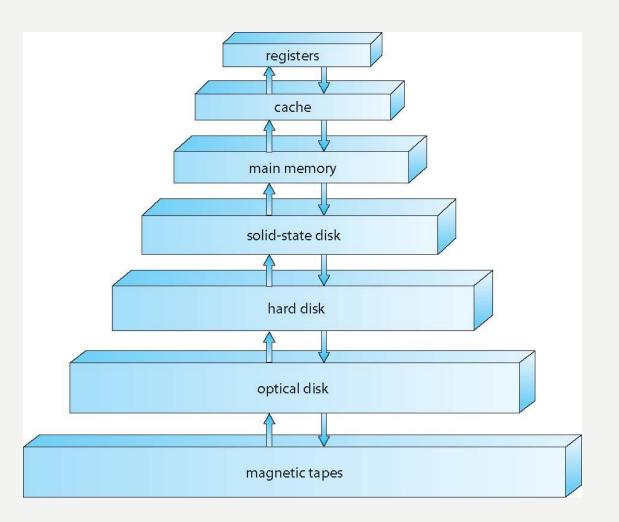
STORAGE STRUCTURE

- Main memory only large storage media that the CPU can access directly
 - Random access
 - Typically **volatile**
- Secondary storage extension of main memory that provides large **nonvolatile** storage capacity
- Hard disks rigid metal or glass platters covered with magnetic recording material
 - Disk surface is logically divided into **tracks**, which are subdivided into **sectors**
 - The disk controller determines the logical interaction between the device and the computer
- Solid-state disks faster than hard disks, nonvolatile
 - Various technologies
 - Becoming more popular

STORAGE HIERARCHY

- Storage systems organized in hierarchy
 - Speed
 - Cost
 - Volatility
- Caching copying information into faster storage system; main memory can be viewed as a cache for secondary storage
- **Device Driver** for each device controller to manage I/O
 - Provides uniform interface between controller and kernel

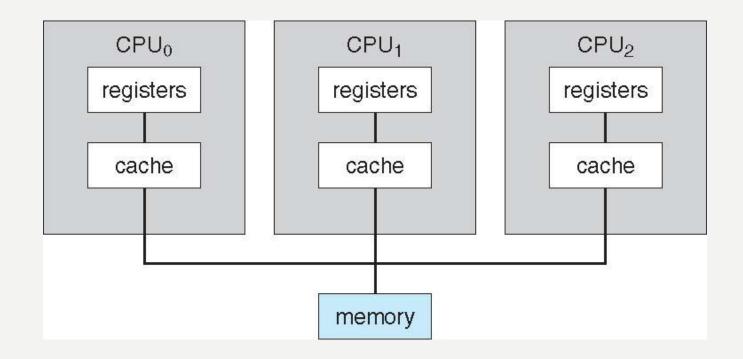
STORAGE-DEVICE HIERARCHY



COMPUTER-SYSTEM ARCHITECTURE

- Most systems use a single general-purpose processor
 - Most systems have special-purpose processors as well
- Multiprocessors systems growing in use and importance
 - Also known as **parallel systems**, **tightly-coupled systems**
 - Advantages include:
 - I. Increased throughput
 - 2. Economy of scale
 - **3.** Increased reliability graceful degradation or fault tolerance
 - Two types:
 - I. Asymmetric Multiprocessing each processor is assigned a specie task.
 - 2. Symmetric Multiprocessing each processor performs all tasks

SYMMETRIC MULTIPROCESSING ARCHITECTURE

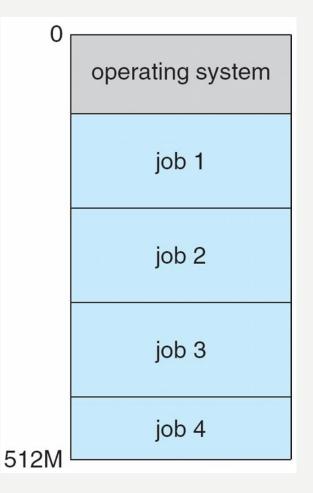


OPERATING SYSTEM STRUCTURE

• Multiprogramming (Batch system) needed for efficiency

- Single user cannot keep CPU and I/O devices busy at all times
- Multiprogramming organizes jobs (code and data) so CPU always has one to execute
- A subset of total jobs in system is kept in memory
- One job selected and run via **job scheduling**
- When it has to wait (for I/O for example), OS switches to another job
- Timesharing (multitasking) is logical extension in which CPU switches jobs so frequently that users can interact with each job while it is running, creating interactive computing
 - Response time should be < I second</p>
 - Each user has at least one program executing in memory ⇒ process
 - If several jobs ready to run at the same time ⇒ CPU scheduling
 - If processes don't fit in memory, **swapping** moves them in and out to run
 - Virtual memory allows execution of processes not completely in memory

MEMORY LAYOUT FOR MULTIPROGRAMMED SYSTEM



MEMORY MANAGEMENT

- To execute a program all (or part) of the instructions must be in memory
- All (or part) of the data that is needed by the program must be in memory.
- Memory management determines what is in memory and when
 - Optimizing CPU utilization and computer response to users
- Memory management activities
 - Keeping track of which parts of memory are currently being used and by whom
 - Deciding which processes (or parts thereof) and data to move into and out of memory
 - Allocating and deallocating memory space as needed

STORAGE MANAGEMENT

- OS provides uniform, logical view of information storage
 - Abstracts physical properties to logical storage unit file
 - Each medium is controlled by device (i.e., disk drive, tape drive)
 - Varying properties include access speed, capacity, data-transfer rate, access method (sequential or random)
- File-System management
 - Files usually organized into directories
 - Access control on most systems to determine who can access what
 - OS activities include
 - Creating and deleting files and directories
 - Primitives to manipulate files and directories
 - Mapping files onto secondary storage
 - Backup files onto stable (non-volatile) storage media

PROTECTION AND SECURITY

- Protection any mechanism for controlling access of processes or users to resources defined by the OS
- Security defense of the system against internal and external attacks
 - Huge range, including denial-of-service, worms, viruses, identity theft, theft of service
- Systems generally first distinguish among users, to determine who can do what
 - User identities (user IDs, security IDs) include name and associated number, one per user
 - User ID then associated with all files, processes of that user to determine access control
 - Group identifier (group ID) allows set of users to be defined and controls managed, then also associated with each process, file
 - Privilege escalation allows user to change to effective ID with more rights